

Improved Perceptual Video Encryption Technique using Residue Number System

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Abstract

In this paper, we propose an enhanced perceptual video encryption technique to speed-up and secure cipher video transmitted across networks using Residue Number System (RNS). The technique proposes a new reverse converter with smaller dynamic range using the moduli set $\{2^n - 1, 2^n, 2^n + 1\}$ that is integrated into our previous work of [2]. After encryption, cipher video is encoded into three residual videos that have smaller pixel values and ideal for transmission across networks. Instead of transmitting the three (3) residual videos, the technique effectively transmits and decodes only two (2) of them back into the original video with same visual quality. Experimental results show that the technique enhances transmission speed and security of cipher video across networks.

1. Introduction

Perceptual video encryption techniques have received overwhelming attention by researchers in recent times leading to the development of varied cryptographic

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advertisement techniques for video of financial interest [3, 4, 7, 11, 13, 15, 30, 31]. These offer content providers the opportunity to reach more viewers while securing the high quality versions of their videos. Potential consumers can preview the degraded videos before acquiring the high quality ones. However, little has been done by previous techniques to enhance the transmission speed of cipher videos transmitted across networks. Owing to large bit representations, big numbers transmit slower than small ones across network. Small numbers can be achieved with the adoption of RNS where residues of the huge and redundant video data are used. Among many inherent features, RNS offers high computational and transmission speeds. In this proposed technique, RNS is employed to split cipher videos in to three shares (residual videos) that have smaller values. Two of the shares are then transmitted and deciphered at the receiver's end without loss of any information. The shares add extra layer of security and enhance transmission speed to the cipher videos. The rest of the paper is presented as follows; residue number system is presented in Section 2. Section 3 presents detail explanation of the proposed reverse converter while Section 4 presents the proposed integrated cryptosystem of the scheme of [2] and RNS. Analysis of results is in Section 5 and finally we conclude in Section 6.

2. Residue Number System

Residue number system is described as a non-weighted number system having numerous benefits in numerical computations. The inherent features in RNS such as the digit-to-digit computations, parallelism, fault tolerance, high computational speed and low power dissipation make it ideal for implementation in fields of communication [9, 10], Digital Signal Processing (DSP) [14, 26, 32], intensive computations such as digital filtering, correlations, convolutions, direct digital frequency synthesis [6], Discrete Fourier Transform (DFT) computations [24], Fast Fourier Transform (FFT) computations, image processing [1, 10, 20, 25, 28] and cryptography [27, 33]. RNS are based on the congruence relation. Two integers x and y are said to be congruent modulo m if m divides exactly the difference of x and y ; mathematically [20]

$$x \equiv y \pmod{m}. \quad (1)$$

If q and r are the quotient and remainder, respectively of the integer division of Y by m that is, $Y = q \cdot m + r$, then, by definition, we have $Y \equiv r \pmod{m}$. The relationship between r ; Y and m is given by [20].

$$r = |Y|_m, \tag{2}$$

RNS is determined by the set $S = m_1, m_2, \dots, m_N$ of N positive and pairwise relatively prime moduli where the dynamic range M is the product of the moduli m_i . Thus;

$$M = \prod_{i=1}^N m_i. \tag{3}$$

With this, every integer Y in $[0, M - 1]$ has a unique representation (y_1, y_2, \dots, y_N) in S . The set S and the numbers y_i are respectively called the moduli set and residue of Y modulo m_i [1, 10, 20].

2.1. RNS to binary conversion

Two traditional techniques are widely used for RNS to binary reverse conversion; the Chinese Remainder Theorem (CRT) and the Mixed Radix Conversion (MRC).

2.1.1. Chinese remainder theorem (CRT)

Consider the moduli set $m_1, m_2, m_3, \dots, m_N$ with $\text{gcd}(m_i, m_j) = 1$ for $i \neq j$ and dynamic range $M = \prod_{i=1}^N m_i$, the CRT for an integer Y having RNS representation $(x_1, x_2, x_3, \dots, x_N)$ is reversed converted as follows;

$$Y = \left| \sum_{i=1}^N M_i |M_i^{-1} x_i|_{m_i} \right|_M, \tag{4}$$

where $M_i = \frac{M}{m_i}$ and M_i^{-1} is the multiplicative inverse of M_i with respect to m_i [17, 20, 29].

2.1.2. The mixed radix conversion (MRC)

Suppose the moduli compose of positive pairwise relative prime integers is being ordered as $m_N > m_{(N-1)} > \dots > m_1$. A non-negative operand Y in the range of $[0, M - 1]$ in mixed radix representation is given by;

$$Y \leftrightarrow \langle \hat{y}_N, \hat{y}_{N-1}, \dots, \hat{y}_1 \rangle,$$

where;

$$Y = \hat{y}_N \prod_{i=1}^{N-1} m_i + \hat{y}_{N-1} \prod_{i=1}^{N-2} m_i + \cdots + \hat{y}_2 m_1 + \hat{y}_1 \quad (5)$$

and $0 \leq \hat{y}_i < m_i$ for all i [17, 19, 20]. Apart from the traditional techniques, other researchers have proposed much faster and efficient RNS to binary reverse converter for particular moduli sets [5, 8, 12, 16, 18, 34]. However, most of them are simple modification of the CRT or MRC techniques [21-23].

3. Proposed Technique

In this section, a new reverse converter with smaller dynamic range is proposed using the moduli set $\{2^n - 1, 2^n, 2^n + 1\}$ for fast decoding of cipher video. The quotient technique is used to convert numbers in RNS to their equivalent binary (weighted) number system. The dynamic range adopted for this technique is $[0, m_3 m_2 - 1]$. Thus, two moduli in the moduli set are used in the reverse conversion process while one modulus is made redundant.

3.1. Proposed reverse converter for dynamic range $[0, \hat{M} = m_3 m_2 - 1]$

Given the pairwise relative prime moduli set $\{m_1, m_2, m_3\} = \{2^n - 1, 2^n, 2^n + 1\}$ and residue (x_1, x_2, x_3) , X can be represented as in the following three different ways;

$$X = m_1 q_1 + x_1, \quad (6)$$

$$X = m_2 q_2 + x_2, \quad (7)$$

$$X = m_3 q_3 + x_3, \quad (8)$$

where $q_i, i = 1, 2, 3$ is the quotient when X is divided by the modulus $m_i, i = 1, 2, 3$. Also, q_2 can be represented as;

$$q_2 = q_3 + d, \quad (9)$$

where $d \in Z$ and $d \geq 0$.

Illustration 1.

Consider $n = 3$, $\{m_1, m_2, m_3\} = \{7, 8, 9\}$, $\hat{M} = 72$ and using Equations (7) and (8)

For $X = 17$,

$$17 = 8(2) + 1 = 9(1) + 8$$

$$\Rightarrow 2 = 1 + 1, \text{ where } d = 1$$

For $X = 32$,

$$32 = 8(4) + 0 = 9(3) + 5$$

$$\Rightarrow 4 = 3 + 1, \text{ where } d = 1$$

For $X = 63$,

$$63 = 8(7) + 7 = 9(7) + 0$$

$$\Rightarrow 7 = 7 + 0, \text{ where } d = 0.$$

Hence, $q_2 = q_3 + d$.

Thus Equation (8) becomes;

$$X = m_2q_3 + m_2d + x_2. \tag{10}$$

From Equation (9);

$$q_3 = \frac{(X - x_3)}{m_3}$$

and substituting gives;

$$X = m_2 \left(\frac{(X - x_3)}{m_3} \right) + m_2d + x_2$$

$$Xm_3 = Xm_2 - m_2x_3 + m_3m_2d + m_3x_2,$$

$$X = \frac{m_3m_2d + m_3x_2 - m_2x_3}{m_3 - m_2}$$

but $(m_3 - m_2) = 1$, hence;

$$X = m_3m_2d + m_3x_2 - m_2x_3.$$

Eliminate the term in d by taking both sides modulo \hat{M}

$$\begin{aligned} |X|_{\hat{M}} &= |m_3m_2d + m_3x_2 - m_2x_3|_{\hat{M}} \\ X &= |m_3x_2 - m_2x_3|_{\hat{M}}. \end{aligned} \tag{11}$$

3.2. Hardware implementation

The binary representations of the respective residues are as follows:

$$x_{1, (n-1)} x_{1, (n-2)} \cdots x_{1, 1} x_{1, 0}, \tag{12}$$

$$x_{2, (n-1)} x_{2, (n-2)} \cdots x_{2, 1} x_{2, 0}, \tag{13}$$

$$x_{3, n} x_{3, n-1} \cdots x_{3, 1} x_{3, 0}. \tag{14}$$

Equation (11) can be simplified as

$$\begin{aligned} X &= |\tau_1 + \tau_2|_{2^n(2^n+1)} \\ &= |\tau_{1, (2n-1)} \tau_{1, (2n-2)} \cdots \tau_{1, 1} \tau_{1, 0} + \tau_{2, (2n-1)} \tau_{2, (2n-2)} \cdots \tau_{2, 1} \tau_{2, 0}|_{2^n(2^n+1)} \\ &= X_{2n-1} X_{2n-2} \cdots X_1 X_0, \end{aligned} \tag{15}$$

where

$$\begin{aligned} \tau_1 &= 2^n x_2 + x_2 \\ &= x_{2, (n-1)} x_{2, (n-2)} \cdots x_{2, 1} x_{2, 0} \overbrace{00 \cdots 00}^{(n\text{-bits})} + x_{2, (n-1)} x_{2, (n-2)} \cdots x_{2, 1} x_{2, 0} \\ &= x_{2, (n-1)} x_{2, (n-2)} \cdots x_{2, 1} x_{2, 0} \boxtimes x_{2, (n-1)} x_{2, (n-2)} \cdots x_{2, 1} x_{2, 0} \\ &= \tau_{1, (2n-1)} \tau_{1, (2n-2)} \cdots \tau_{1, 1} \tau_{1, 0} \end{aligned} \tag{16}$$

and,

$$\begin{aligned} \tau_2 &= |-2^n x_3|_{2^n(2^n+1)} \\ &= |\bar{x}_{3, n} \bar{x}_{3, (n-1)} \cdots \bar{x}_{3, 1} \bar{x}_{3, 0} \overbrace{(11 \cdots 11)}^{(n\text{-bits})}|_{2^n(2^n+1)} \\ &= \tau_{2, (2n-1)} \tau_{2, (2n-2)} \cdots \tau_{2, 1} \tau_{2, 0} \end{aligned} \tag{17}$$

3.3. Hardware realization

The hardware of the proposed scheme applicable to values of $n \leq 4$, can be realized with a simple modulo carry propagate adder (MCPA) and an inverter. Equation (16) is a concatenation of bits, which would not require any hardware resource for that operation. Equation (17) would only require a bit inverter, which is not expensive and at the same time, does not impose undue delay on the scheme. It is only in Equation (15) that an MCPA of length $2n$ -bits would be required to add the results of (16) and (17). Thus, the hardware resources in this regard are $2n$ -bits wide while the delay imposed by such an adder is $4n$ -bits. The delay for an inverter is usually unity; therefore, the total delay that would be imposed on the scheme is $(4n + 1)$ -bits. The block diagram for the reverse conversion is shown in Figure 1.

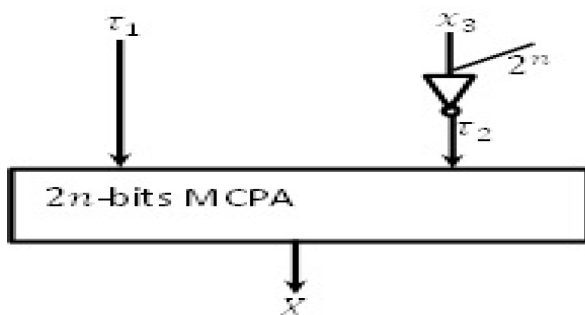


Figure 1. Block diagram for the reverse conversion.

4. RNS Integration with Proposed Video Cryptosystem

4.1. Enhanced encryption algorithm

The proposed enhancement to the encryption algorithm of [2] is obtained by modifying the last step to include the RNS encoder. This is presented as;

- Input video file V , number of iteration I , block size $blkSize$, angle of rotation θ , unit-anti-diagonal matrix K and the number of bits n .
- Encrypt video frames as follows:
 - Compute point of rotation and extract block of pixels (V_{P_b}) using rotation matrix (A) .

- Multiply block of pixels (V_b) by unit anti-diagonal matrix (K) to obtain cipher frame ($V_{be} = V_{P_b} \times K$).
- Add cipher frame to cipher video ($V_{E+} = V_{be}$).
- Compute the moduli set $\{2^n - 1, 2^n, 2^n + 1\}$ using n .
- Compute the residual videos x_2 and x_3 from V_E modulus 2^n and $2^n + 1$, respectively.
- Transmit cipher videos x_2 and x_3 .

4.2. Enhanced decryption algorithm

The proposed enhancement to the decryption algorithm of [2] is obtained by modifying the initial steps to include the RNS decoder. The proposed enhancement is presented as;

- Input cipher videos x_2 and x_3 , number of iteration I , block size $blkSize$, angle of rotation θ , anti-diagonal matrix K and the number of bits n .
- Compute the moduli set $\{2^n - 1, 2^n, 2^n + 1\}$ using n .
- Decrypt cipher frames as follows:
 - Decode x_2 and x_3 back into cipher frames V_E using equation (11).
 - Compute point of rotation and extract block of pixels V_{E_b} using rotation matrix (A).
 - Multiply block of pixels V_{E_b} by unit anti-diagonal matrix (K) to obtain plain frame $V_{bp} = V_{E_b} \times K$.
 - Add plain frame to plain video $V+ = V_{bp}$.
- Transmit cipher video V .

In Figure 2, the forward and proposed reverse converters for the moduli set $\{2^n - 1, 2^n, 2^n + 1\}$ are fitted after the encryption algorithm and before the decryption algorithm respectively. Pixels values of encrypted video are passed through the forward

converter which yields two residues (x_2, x_3) corresponding to the moduli $\{2^n\}$ and $\{2^n + 1\}$. Continuous bitstream of the residues are transmitted through the transmission channel in fixed length code words. At the receivers' end, the fixed length code words are transformed back to continuous form and the proposed reverse converter is applied to recover the cipher video. The original video is then recovered from the cipher video by applying the decryption algorithm sub-block.

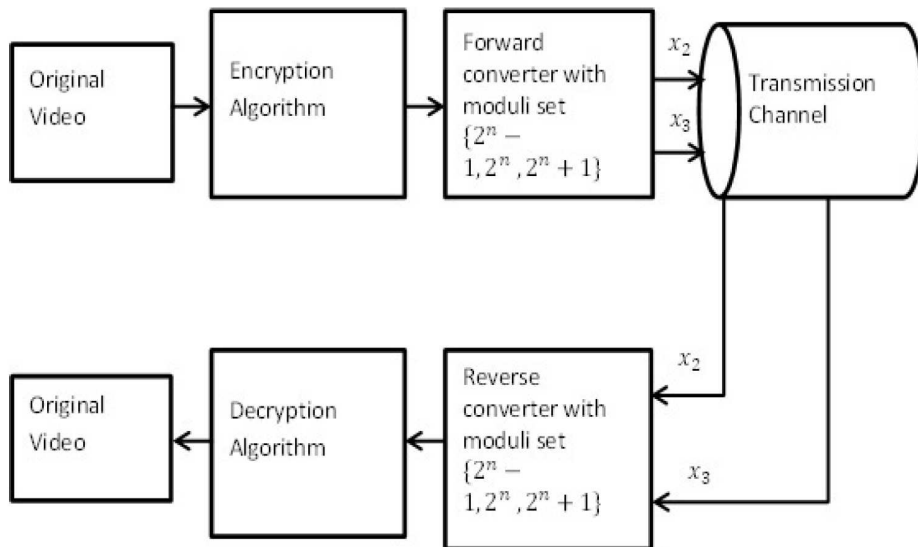


Figure 2. Block diagram of proposed integration of RNS with perceptual video cryptosystems.

The MATLAB Simulink system to test the proposed integrated scheme is shown in Figure 3. RNS is integrated with the work of [2]. After encryption with ‘Encryption Algorithm1’ sub-block, the cipher video is passed through the sub-block ‘RNSEncoder’ to be encoded into two residual videos x_2 and x_3 for transmission. Since residual values are less than that of the originals, transmission is much easier and faster. At the receiver’s end, residual videos are decoded by the ‘RNSDecoder’ sub-block into the cipher video followed by decryption into the plain video using the ‘Decryption Algorithm1’ sub-block.

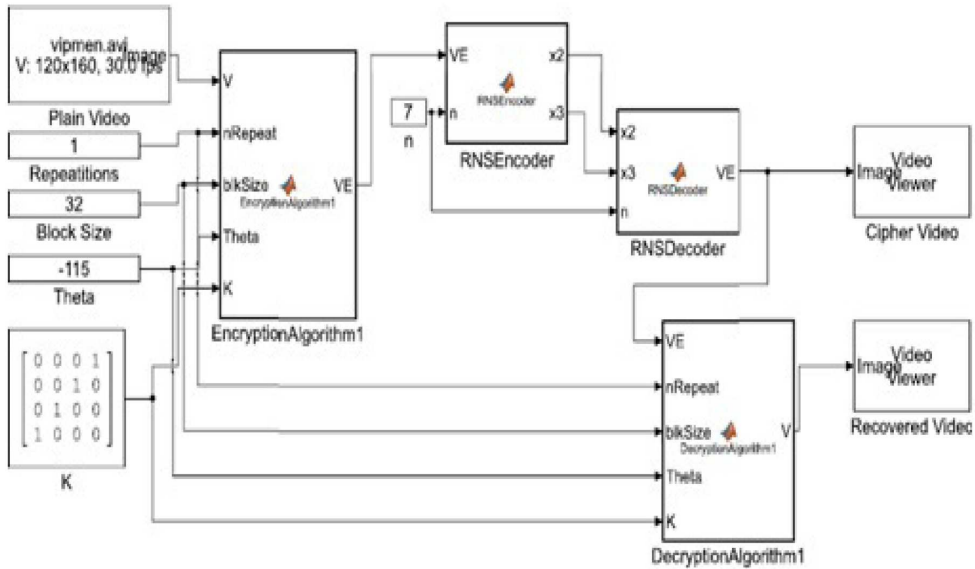


Figure 3. Integration of RNS with proposed perceptual video via unit anti-diagonal matrix.

Table 1. Processing time of encryption/decryption algorithm.

Video File(Dimension)	Processing Time(s)	
	Encryption(RNSEncoder)	Decryption(RNSDecoder)
vipmen.avi (160 × 120)	0.058(0.008)	0.061(0.026)
carphone.avi (176 × 144)	0.092(0.014)	0.098(0.035)
xylophone.mpg (320 × 240)	0.468(0.036)	0.475(0.097)

Figure 4 shows experimental results of the residual videos after encoding a cipher frame using $n = 4$. The dark pictures of parts (b) and (c) confirm the smaller values achieved through the RNS encoder. Any adversary receiving these will need extra efforts to decode them to the cipher state before efficient decryption can occur.

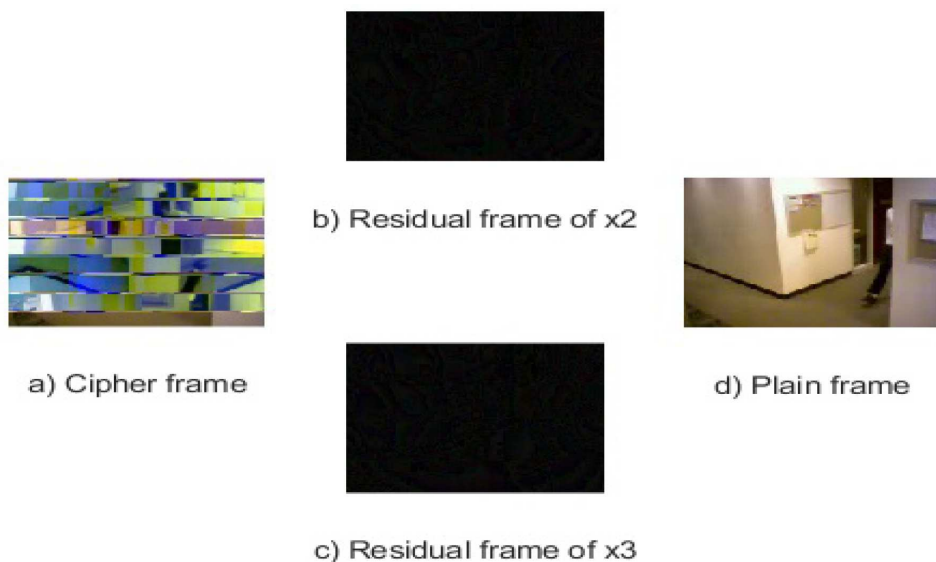


Figure 4. Residual frames of encoded video using RNS. (a) Cipher frame, (b) Residual frame x_2 , (c) Residual frame x_3 and (d) Recovered frame.

5. Analysis of Results

5.1. Procession time

The average processing time (in seconds) for ten (10) simulations of the encryption and decryption processes are summaries in Table 1. The results indicate that the 'RNSEncoder' consumes about 9% of the average encryption time while 25% of the average decryption time is consumed by the 'RNSDecoder'. Thus, RNS constitutes about 34% of the total processing time of any given encryption and decryption operation. Also, the increase in the decryption time enhances security especially against Bruce-force

5.2. Encoding analysis

Video data is made up of several frames(images) arranged over time. The elements of these frames are called pixels. The weight of a pixel for the Red-Green-Blue (RGB) format of video span 0 to 2^8 . Thus, 8-bits are required to encode and transmit each pixel. However, the introduction of RNS reduces this to 4-bits when $n = 4$ is used for the 'RNSEncoder'. This saves half of the bits required to encode and transmit cipher videos. Consequently, transmission speed is enhanced by the introduction of RNS.

6. Conclusion

This paper proposes an enhancement to the perceptual video encryption and decryption algorithms proposed in [2] using RNS. Analysis of simulated results shows that the enhanced scheme increases transmission speed and adds extra security to cipher video.

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